## Low Data Complexity Attacks on AES

#### Orr Dunkelman

#### Joint work with Nathan Keller Faculty of Mathematics and Computer Science Weizmann Institute of Science

#### 9 February, 2010



Results

## A Cryptographic Challenge

#### Consider one full round of AES:



Chal	lenge

#### Results

# Why?

- A LEX variant (one that leak the entire state every 4 rounds).
- A sub-routine in other attacks (in case of KP attacks).
- ALPHA MAC?
- ▶ Hash functions? (Vortex-0.0)

Results

# Why?

- A LEX variant (one that leak the entire state every 4 rounds).
- A sub-routine in other attacks (in case of KP attacks).
- ALPHA MAC?
- ▶ Hash functions? (Vortex-0.0)
- Practicing the art of running after bits.

#### So How Far Can You Go with 1 KP on AES-128

Seems easy ah?

## So How Far Can You Go with 1 KP on AES-128

- Seems easy ah?
- Our best attack takes 2<sup>48</sup> using (heavily) the key schedule.
- We asked some people doing algebraic attacks, and they could not do better.
- We have strong evidence to believe that it is impossible to go below 2<sup>32</sup>.

#### Results

### How Far Can You Go With 1 KP on AES-128

- Consider 2 full rounds of AES with 1 KP.
- Not so easy, ah?

#### Results

## How Far Can You Go With 1 KP on AES-128

- Consider 2 full rounds of AES with 1 KP.
- Not so easy, ah?
- Our best attack takes 2<sup>96</sup> using (heavily) the key schedule.

Results

## Some More Results

Attack Type	Number of	Complexity		Independent
	Rounds	Data	Time	Keys
MitM	1	1 KP	2 <sup>48</sup>	
Diff.	1	2 KP	2 <sup>16</sup>	
MitM	2	1 KP	2 <sup>96</sup>	
Diff.	2	2 KP	2 <sup>48</sup>	
Diff.	2	3 KP	2 <sup>32</sup>	
Diff.	3	2 CP	2 <sup>28</sup>	$\checkmark$
Diff.	3	9 KP	2 <sup>40</sup>	$\checkmark$
Diff.	4	5 CP	2 <sup>64</sup>	$\checkmark$
Diff. MitM	4	11 CP	2 <sup>40</sup>	$\checkmark$
Diff. MitM	4	2 <sup>49.3</sup> KP	2 <sup>49.3</sup>	$\checkmark$
Diff. MitM	4	2 <sup>33.2</sup> KP	2 <sup>80</sup>	$\checkmark$
Diff. MitM	4	2 <sup>17.7</sup> KP	2 <sup>120</sup>	$\checkmark$

Results

## Side Effects

#### AES Proposal: Rijndael

"In order to make the cipher and its inverse more similar in structure, the linear mixing layer of the last round is different from the mixing layer in the other rounds. It can be shown that this does not improve or reduce the security of the cipher in any way. This is similar to the absence of the swap operation in the last round of the DES." Challenge Results
Side Effects (cont.)

- If the last rounds MixColumns is dropped, our 1-round attack with 1 KP is improved to 2<sup>16</sup>!
- Other attacks improve as well.
- Hence, the lack of MixColumns in the last round does affect security.
- Note that breaking the symmetry may help against slide attacks.

For more information: See http://eprint.iacr.org/2010/041.

Results

#### Questions?

## Thank you for your attention.